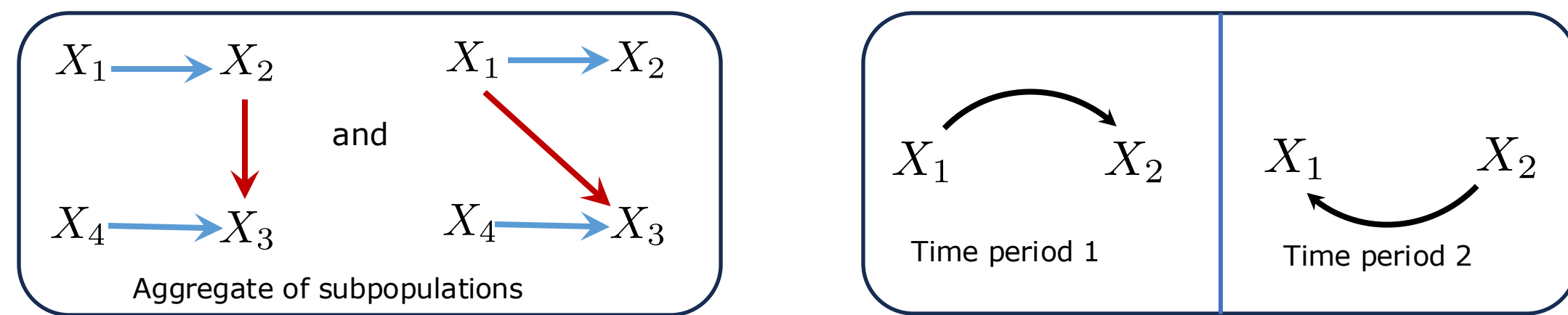




Motivation

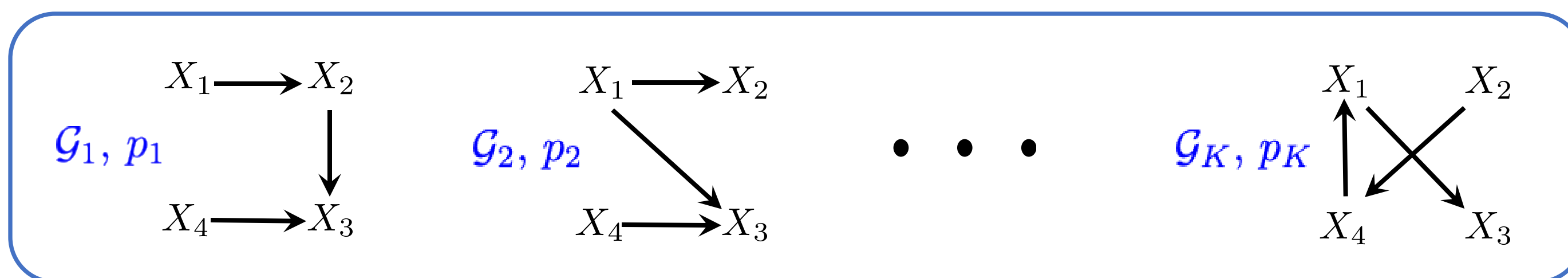
Most real causal systems are complex

- **Time-varying systems:** causal relationships can change over time
- **Cyclic relationships:** e.g., feedback loops
- **Modeling subpopulations:** e.g., subtypes of cancers do not share the same exact biological pathways



These complex models are better modeled by a mixture of DAGs!

Mixture Model



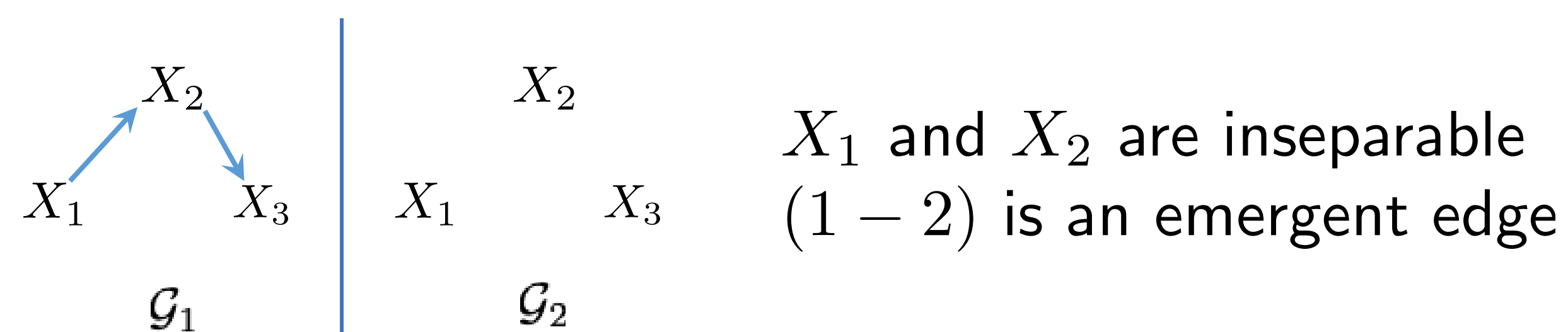
- A mixture of (unknown) K DAGs over n nodes: $\{\mathcal{G}_1, \dots, \mathcal{G}_K\}$
- Mixture distribution:

$$p_{\text{mix}} = \sum_{\ell \in [K]} w_{\ell} \cdot p_{\ell}, \quad \text{where } \sum_{\ell \in [K]} w_{\ell} = 1$$

- **True edges:** exist in at least one component DAG
- Define **mixture parents:** $\text{pa}_{\text{mix}}(i) = \bigcup_{\ell \in [K]} \text{pa}_{\ell}(i)$

Challenges of causal discovery in mixtures

- Single DAG: CI tests on observational data give the skeleton
- Mixture of DAGs: spurious unbreakable dependencies



Prior work on **observational data**

- FCI algorithm, graphical restrictions e.g. no cycles (Saeed et al. 2020)
- Longitudinal data, no solution for spurious (Strobl, 2023)
- Necessary-sufficient conditions for **emergent** edges (Varici et al. TMLR 2024) (*Separability Analysis for Causal Discovery in Mixture of DAGs*)

Observational data is not sufficient

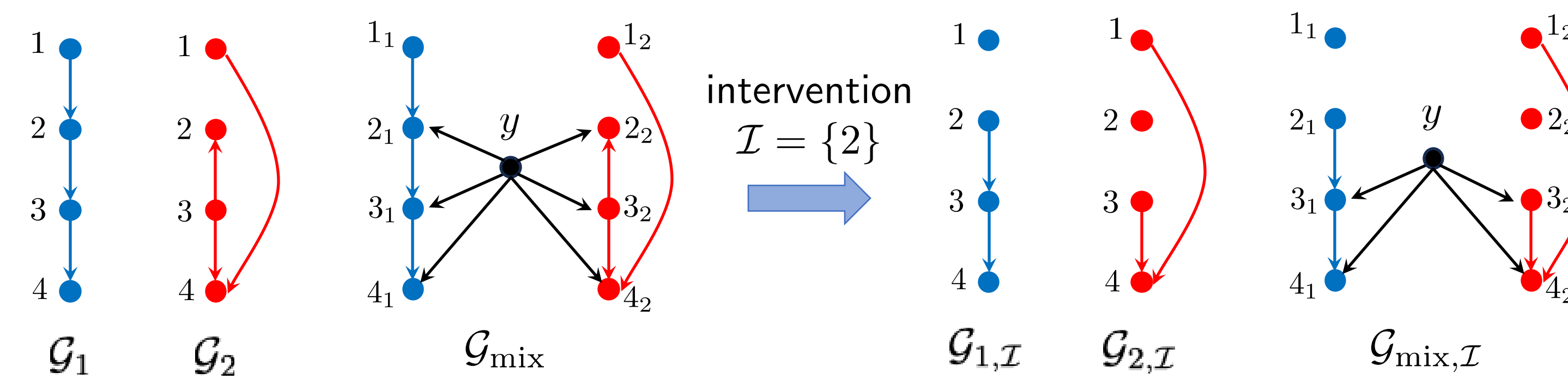
Interventional Causal Discovery

Intervention model: For a subset of nodes $I \subset [n]$, cut off parents

$$p_{\ell}(x_i | x_{\text{pa}_{\ell}(i)}) \rightarrow q_i(x_i), \quad \forall i \in I, \forall \ell \in [K]$$

$$p_{\ell, \mathcal{I}}(x) \triangleq \prod_{i \in \mathcal{I}} q_i(x_i) \prod_{i \in \mathbf{V} \setminus \mathcal{I}} p_{\ell}(x_i | x_{\text{pa}_{\ell}(i)}), \quad \forall \ell \in [K]$$

interventional mixture dist.: $p_{\text{mix}, \mathcal{I}} = \sum_{\ell \in [K]} w_{\ell} \cdot p_{\ell, \mathcal{I}}$



- True edges: $\{(1 \rightarrow 2), (2 \rightarrow 3), (3 \rightarrow 2), (3 \rightarrow 4), (1 \rightarrow 4)\}$
- Mixture parents, e.g., $\text{pa}_{\text{mix}}(2) = \{1, 3\}$ and $\text{pa}_{\text{mix}}(3) = \{2\}$
- Spurious dependencies: $\{(1, 3), (2, 4)\}$

goal: identify true edges (or mixture parents) via interventions

1. Necessary and sufficient sizes of interventions
2. A learning algorithm with near-optimal interventions

Assume interventional mixture faithfulness (standard extension)

Results

Theorem (intervention size): To find $\text{pa}_{\text{mix}}(i)$ of via CI tests

1. interventions with size $|I| \leq |\text{pa}_{\text{mix}}(i)| + 1$ are **sufficient**
2. at the worst-case, $|I| = |\text{pa}_{\text{mix}}(i)| + 1$ is **necessary**

Why? To determine whether $j \in \text{pa}_{\text{mix}}(i)$, $I = \text{pa}_{\text{mix}}(i) \cup \{j\}$ suffices

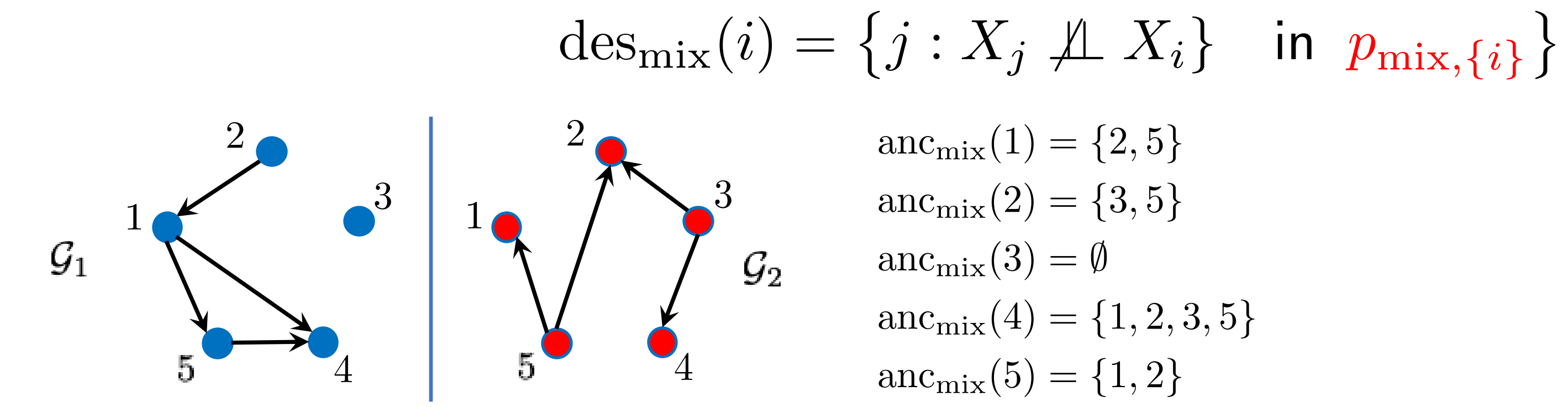
Theorem (int. size - trees): For a mixture of trees

1. interventions with size $|I| \leq K + 1$ are **sufficient**
2. at the worst-case, $|I| = K + 1$ is **necessary**

Why? There are at most K paths from j to i , one for each DAG
At the worst case (all disjoint paths), we need to block them all

Learning Algorithm

Step 1. **mixture ancestors/descendants** via atomic interventions



Repeat steps 2, 3, 4 for all nodes $i \in \mathbf{V}$

Step 2. **Break cycles among ancestors**

- $\mathcal{C}(i)$: all cycles among $\text{anc}_{\text{mix}}(i)$
- $\mathcal{B}(i)$: minimal set that intersects with every cycle in $\mathcal{C}(i)$
e.g. $\mathcal{C}(4) = \{(2, 5, 2), (2, 1, 5, 2), (1, 5, 1)\} \rightarrow \mathcal{B}(4) = \{5\}$
- **Cyclic complexity:** $\tau_i \triangleq |\mathcal{B}(i)|$
- $I = \mathcal{B}(i) \cup \{j\}$ for all $j \in \text{anc}_{\text{mix}}(i)$ to refine descendants

Step 3. **Topological layering:** refined descendant sets do not conflict!

- Bottom-up layering: $S_1(4) = \{1\}$, $S_2(4) = \{2\}$, $S_3(4) = \{3, 5\}$

Step 4. **Identify mixture parents:** process each layer sequentially

- For every possible parent j , intervene on $\mathcal{B}(i) \cup \text{pa}_{\text{mix}}(i) \cup \{j\}$
- e.g. $I = \{1, 5\}$ determines whether $1 \in \text{pa}_{\text{mix}}(4)$

Theorem: Learns all true edges using $\mathcal{O}(n^2)$ interventions, with interv. size at most $|\text{pa}_{\text{mix}}(i)| + \tau_i + 1$ for each node i .

Optimality gap = cyclic complexity of node i

Experiments

- Mixture of Gaussians, vary number of DAGs and graph size
- Strong performance for all settings.
- Empirical average cyclic complexity: less than 2 (for 10 nodes, $K = 3$)

